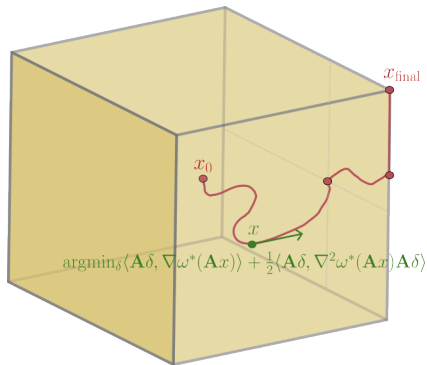


# Discrepancy Minimization via Regularization



Lucas Pesenti



Bocconi University

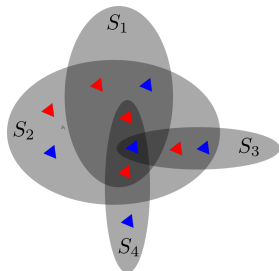
Adrian Vladu



CNRS & Université Paris Cité

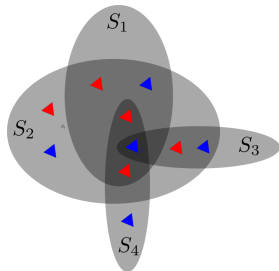
# Discrepancy Theory

**Discrepancy/Vector Balancing problem:** Given  $u_1, \dots, u_n \in K \subseteq \mathbb{R}^d$ , find  $x_1, \dots, x_n \in \{\pm 1\}$  s.t. the *discrepancy*  $\|\sum_i x_i u_i\|$  is “small”



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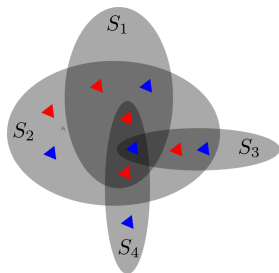


## Examples:

- ▶ (Spencer's theorem)  $K := \{u : \|u\|_\infty \leq 1\}$ , target  $\|\cdot\|_\infty$
- ▶ (Komlós conjecture)  $K := \{u : \|u\|_2 \leq 1\}$ , target  $\|\cdot\|_\infty$
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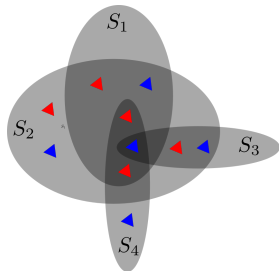
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## Motivations:

- ▶ Prove the existence of rare objects
- ▶ Toy problem/building block for “sparsification” tasks

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*Our contribution:*



- ▶ A new algorithmic framework for these problems
- ▶ A tighter constant in [Spencer's theorem](#)
- ▶ A proof of [Komlós conjecture](#) for “pseudorandom” inputs

# Discrepancy and Continuous Methods

**Thm:** [Spencer'85] For any  $\mathbf{A} \in \mathbb{R}^{n \times n}$  s.t.  $|\mathbf{A}_{ij}| \leq 1$ , there exists  $x \in \{\pm 1\}^n$  s.t.  $\|\mathbf{A}x\|_\infty = O(\sqrt{n})$

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Many different *algorithmic* proofs:

- ▶ [Bansal'10, Lovett-Meka'12] random walks, SDP
- ▶ [Eldan-Singh'14, Rothvoss'14] LP with random objective
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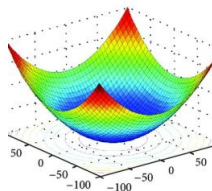
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More and more inspired by “continuous” optimization



Our algorithm: **Newton's method** on a **regularized** objective



# Spencer's Theorem via Regularization (1/3)

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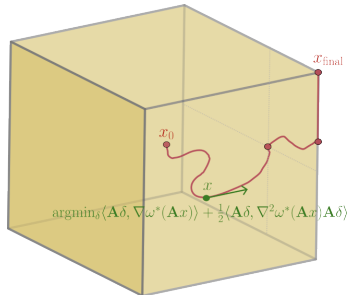
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While  $x \notin \{\pm 1\}^n$

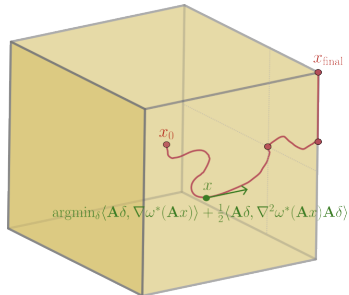
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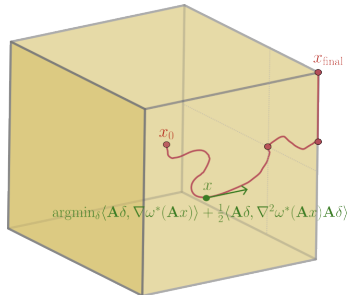
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$x$  makes a small step in direction  $\delta$  while staying in  $[-1, 1]^n$

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**Def:** *Regularized* maximum

$$\omega^*(y) := \max_{r \in \Delta_n} \langle y, r \rangle + \sum_{i=1}^n r_i^{\frac{1}{2}}.$$

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**Claim:**  $\omega^*(\mathbf{A}x) = \|\mathbf{A}x\|_\infty \pm O(\sqrt{n})$



# Spencer's Theorem via Regularization (3/3)

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Analysis idea:

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Analysis idea:

1. By studying  $\nabla^2 \omega^*(\mathbf{A}x)$ , we prove that there always exists a direction  $\delta \perp x$ ,  $\text{supp}(\delta) \subseteq F$  s.t.

$$\omega^*(\mathbf{A}(x + \delta)) - \omega^*(\mathbf{A}x) \leq \frac{\|\delta\|_2^2}{\sqrt{|F|}}.$$

2. Hence we get charged  $\frac{1}{\sqrt{|F|}}$  cost “per unit of  $\|x\|_2^2$ ”
3. Worst case: coordinates get frozen every time  $\|x\|_2^2$  increases by 1, total cost

$$1 \times \frac{1}{\sqrt{n}} + 1 \times \frac{1}{\sqrt{n-1}} + \dots = O(\sqrt{n}). \quad \square$$

# Our Results (1/2)

Improved Constant for Spencer's Theorem:

**Thm:** [P-V'23] For any  $\mathbf{A} \in \mathbb{R}^{n \times n}$  s.t.  $|\mathbf{A}_{ij}| \leq 1$ , there is  $x \in \{\pm 1\}^n$   
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Proof idea: slightly different regularizer, track **constants** carefully

## Our Results (2/2)

Komlós conjecture for “pseudorandom” inputs:

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where  $\lambda := \max_{\|v\|_2=1, v \perp \mathbf{1}} \|\mathbf{A}^{\odot 2}v\|_2$  ( $\odot \equiv$  entrywise product)

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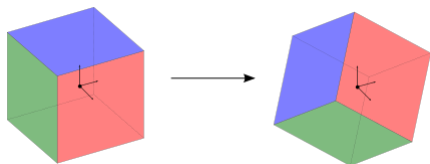
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- ▶ Generalizes Banaszczyk's bound and [Potukuchi'18]
- ▶ Special case:



“A randomly rotated hypercube has a corner at  $\infty$ -distance  $O(1)$  from the origin”

*Open problem:* Komlós conjecture for worst-case rotations?

# Conclusion

*Our algorithm:* **Newton's method** on a **regularized** objective

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Future directions:

- ▶ Prove a tight constant in Spencer's theorem
- ▶ Application of this framework to matrix discrepancy and graph sparsification
- ▶ Bridge this framework and arguments based on Lovász local lemma ( $\rightsquigarrow$  Beck-Fiala conjecture)